

Multi-Armed Bandits with Relative Feedback

Pratik Gajane

Orange labs & INRIA SequeL

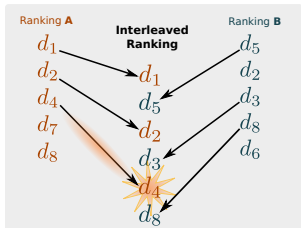
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Outline

1. Dueling bandits
2. Analysis of the algorithm
3. Experiments

Motivation for the dueling bandit problem

- In many practical situations, **relative feedback** is available, and not absolute feedback.
- Eg: “I like Tennis more than Basketball” instead of “I value tennis at 48/50 and Basketball at 33/50”.



- Information retrieval systems where users provide *implicit feedback* about the provided results.
- **Interleaved filtering**, proposed by Radlinski et al. [3], interleaves the rankings to remove the bias.
- Inability of classical MAB to deal with **relative feedback** motivates new problem setting.

The dueling bandit problem

- A variation of the classical Multi-Armed Bandit (MAB) to deal with **relative feedback**.
- At each time period, the learner selects two arms.
- The learner only sees the outcome of the *duel* between the selected arms.
- The learner receives a function of the rewards of the selected arms.



Formulating the dueling bandits

- Matrix-based formulation

- ▶ *preference matrix* contains $\mathbb{P}_{a,b}$ = unknown probability with which a wins the duel against b .

$$\begin{array}{c} \\ 1 \\ 2 \\ \vdots \\ K \end{array} \begin{bmatrix} & 1 & 2 & \dots & K \\ 1 & 1/2 & \mathbb{P}_{1,2} & & \mathbb{P}_{1,K} \\ 2 & \mathbb{P}_{2,1} & 1/2 & & \mathbb{P}_{2,K} \\ & & & \ddots & \\ K & \mathbb{P}_{K,1} & \mathbb{P}_{K,2} & & 1/2 \end{bmatrix}$$

- Utility-based formulation

- ▶ At each time t , a utility $x_a(t)$ is associated with each arm a .

- ▶ When arms a and b are selected,

$x_a(t) > x_b(t)$: a wins the duel

$x_a(t) < x_b(t)$: b wins the duel

$x_a(t) = x_b(t)$: $\begin{cases} a \text{ wins the duel with probability } 0.5 \\ b \text{ wins the duel with probability } 0.5 \end{cases}$

Utility-based adversarial dueling bandits

- State of the art dueling bandits algorithms are for stochastic bandits. \rightarrow arm rewards are **independent and identically distributed (iid)**.
- **Adversarial** dueling bandits allow us to drop these assumptions.
- In our setting, the adversary chooses a sequence of utility vectors $\mathbf{x}(t) = (x_1(t), \dots, x_K(t)) \in [0, 1]^K$ for $t = 1, \dots, T$.
- At each time t , the learner chooses two arms a and b ,
Instantaneous reward = $\frac{x_a(t) + x_b(t)}{2}$ (hidden)

$$\mathbf{Feedback} = x_a(t) - x_b(t)$$

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- In our setting, the adversary chooses a sequence of utility vectors $\mathbf{x}(t) = (x_1(t), \dots, x_K(t)) \in \{0, 1\}^K$ for $t = 1, \dots, T$.
- At each time t , the learner chooses two arms a and b ,

$$\begin{aligned} \text{Instantaneous reward} &= \frac{x_a(t) + x_b(t)}{2} \quad (\text{hidden}) \\ \text{Feedback (binary rewards)} &= \begin{cases} -1 & \text{if } x_a(t) < x_b(t) \\ 0 & \text{if } x_a(t) = x_b(t) \\ +1 & \text{if } x_a(t) > x_b(t) \end{cases} \end{aligned}$$

Lower bound for any dueling bandit algorithm

Theorem

For $K \geq 2$ and $T \geq K$, there exists a distribution over assignments of rewards such that the **expected cumulative regret** of any utility-based dueling bandit algorithm cannot be less than $\Omega(\sqrt{KT})$.

\mathbb{G}_{max} - Maximum possible reward for a single-arm strategy

$\mathbb{E}(\mathbb{G}_{alg})$ - Expected reward earned by the algorithm's strategy

$\mathbb{G}_{max} - \mathbb{E}(\mathbb{G}_{alg})$ - **Expected cumulative regret**

- We proved this by reduction to classical bandits as suggested in Ailon et al. [1]
- Lower bound for adversarial dueling bandits = lower bound of classical adversarial bandits = $\Omega(\sqrt{KT})$
- Data dependent lower bound for stochastic bandits = $\Omega(K \log(T)/\Delta)$

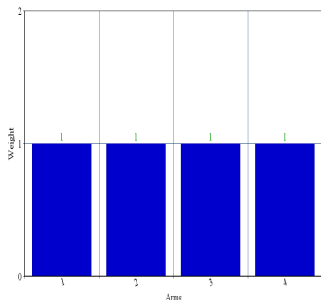
Relative Exponential Weighing Algorithm (REX3)

- Non-trivial extension of EXP3 [2] to the dueling bandits with binary rewards.
- Assigns a weight to each arm. Higher weight \implies higher selection probability.
- $d = x_a - x_b = \begin{cases} -1 & \text{if } x_a < x_b \\ 0 & \text{if } x_a = x_b \\ +1 & \text{if } x_a > x_b \end{cases}$
- For anytime version, a kind of “doubling trick” (Seldin et al. [4]).

- 1: **Parameters:** Real $\gamma \in (0, 0.5)$
- 2: **Initialization:** $w_i(1) = 1$ for $i = 1, \dots, K$.
- 3: **for** $t = 1, 2, \dots$ **do**
- 4: **for** $i = 1, \dots, K$ **do**
- 5: $p_i(t) \leftarrow (1 - \gamma) \frac{w_i(t)}{\sum_{j=1}^K w_j(t)} + \frac{\gamma}{K}$
- 6: **end for**
- 7: Pull $a, b \sim (p_1(t), \dots, p_K(t))$.
- 8: Get relative feedback $d \in \{-1, 0, +1\}$
- 9: **if** $a \neq b$ **then**
- 10: $w_a(t+1) \leftarrow w_a(t) \cdot e^{\frac{\gamma}{K} \frac{d}{2p_a}}$
- 11: $w_b(t+1) \leftarrow w_b(t) \cdot e^{-\frac{\gamma}{K} \frac{d}{2p_b}}$
- 12: **end if**
- 13: Update γ (for anytime version)

Relative Exponential Weighing Algorithm (REX3)

Weights at $t = 0$
($\gamma = 0.4$)

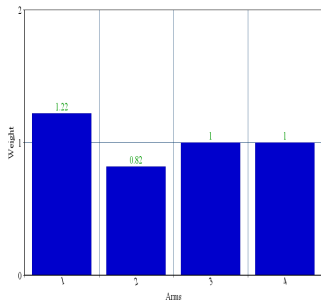


- Update weight according to (relative) feedback.

- 1: **Parameters:** Real $\gamma \in (0, 0.5)$
- 2: **Initialization:** $w_i(1) = 1$ for $i = 1, \dots, K$.
- 3: **for** $t = 1, 2, \dots$ **do**
- 4: **for** $i = 1, \dots, K$ **do**
- 5: $p_i(t) \leftarrow$
 $(1 - \gamma) \frac{w_i(t)}{\sum_{j=1}^K w_j(t)} + \frac{\gamma}{K}$
- 6: **end for**
- 7: Pull
 $a, b \sim (p_1(t), \dots, p_K(t))$.
- 8: Get relative feedback
 $d \in \{-1, 0, +1\}$
- 9: **if** $a \neq b$ **then**
- 10: $w_a(t+1) \leftarrow w_a(t) \cdot e^{\frac{\gamma}{K} \frac{d}{2p_a}}$
- 11: $w_b(t+1) \leftarrow w_b(t) \cdot e^{-\frac{\gamma}{K} \frac{d}{2p_b}}$
- 12: **end if**
- 13: Update γ (for anytime version)

Relative Exponential Weighing Algorithm (REX3)

$a = 1, b = 2, x_a > x_b$
Weights at $t = 1$

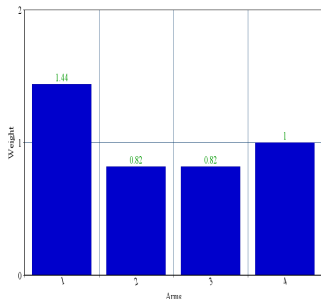


- Weight may decrease unlike EXP3.

- 1: **Parameters:** Real $\gamma \in (0, 0.5)$
- 2: **Initialization:** $w_i(1) = 1$ for $i = 1, \dots, K$.
- 3: **for** $t = 1, 2, \dots$ **do**
- 4: **for** $i = 1, \dots, K$ **do**
- 5: $p_i(t) \leftarrow$
 $(1 - \gamma) \frac{w_i(t)}{\sum_{j=1}^K w_j(t)} + \frac{\gamma}{K}$
- 6: **end for**
- 7: Pull
- 8: $a, b \sim (p_1(t), \dots, p_K(t))$.
- 9: Get relative feedback
 $d \in \{-1, 0, +1\}$
- 10: **if** $a \neq b$ **then**
- 11: $w_a(t+1) \leftarrow w_a(t) \cdot e^{\frac{\gamma}{K} \frac{d}{2p_a}}$
- 12: $w_b(t+1) \leftarrow w_b(t) \cdot e^{-\frac{\gamma}{K} \frac{d}{2p_b}}$
- 13: **end if**
- 14: Update γ (for anytime version)

Relative Exponential Weighing Algorithm (REX3)

$a = 1, b = 3, x_a > x_b$
Weights at $t = 2$



- Weights spike at arms who win the duel regularly.

- 1: **Parameters:** Real $\gamma \in (0, 0.5)$
 - 2: **Initialization:** $w_i(1) = 1$ for $i = 1, \dots, K$.
 - 3: **for** $t = 1, 2, \dots$ **do**
 - 4: **for** $i = 1, \dots, K$ **do**
 - 5: $p_i(t) \leftarrow (1 - \gamma) \frac{w_i(t)}{\sum_{j=1}^K w_j(t)} + \frac{\gamma}{K}$
 - 6: **end for**
 - 7: Pull
 $a, b \sim (p_1(t), \dots, p_K(t))$.
 - 8: Get relative feedback
 $d \in \{-1, 0, +1\}$
 - 9: **if** $a \neq b$ **then**
 - 10: $w_a(t+1) \leftarrow w_a(t) \cdot e^{\frac{\gamma}{K} \frac{d}{2p_a}}$
 - 11: $w_b(t+1) \leftarrow w_b(t) \cdot e^{-\frac{\gamma}{K} \frac{d}{2p_b}}$
 - 12: **end if**
 - 13: Update γ (for anytime version)
- 9

Upper bound for REX3

Theorem

$$G_{max} - \mathbb{E}(G_{alg}) \leq \frac{K}{\gamma} \ln(K) + \gamma T$$

where

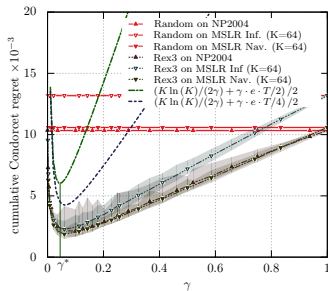
$$\tau = e \cdot \mathbb{E}G_{alg} - (4 - e) \cdot \mathbb{E}G_{unif}$$

Corollary

When $\gamma = \min \left\{ \frac{1}{2}, \sqrt{\frac{K \ln(K)}{\tau}} \right\}$, the *expected cumulative regret* of REX3 is bounded by $\mathcal{O} \left(\sqrt{K \ln(K) T} \right)$.

- Upper bound of REX3 = Upper bound of EXP3.
- Optimality:

$$\text{REX3} \sim_{\ln} \Omega \left(\sqrt{KT} \right)$$



Analysis of REX3

- Main challenge of dueling bandits: no direct way to estimate absolute reward values like EXP3.
- In EXP3, since we can observe absolute feedback (x_a), the estimator $\hat{x}_i(t)$ is defined as follows:

$$\hat{x}_i(t) = \mathbb{I}[i = a] \frac{x_a(t)}{p_a(t)}$$

- The division by p_a ensures that more “surprising” (i.e. lower p_a) the observed reward x_a , higher is the estimator.
- Ensures that their expectations are equal to the actual rewards for each action i.e.

$$\mathbb{E}[\hat{x}_i(t)] = x_i(t)$$

Analysis of REX3

- Feedback in dueling bandits is relative ($x_a - x_b$) instead of absolute (x_a), so the use of EXP3 estimator is not possible.
- To overcome this challenge, we introduced a new estimator $\hat{c}_i(t)$.
- We define $\hat{c}_i(t)$ in the following way:

$$\hat{c}_i(t) = \mathbb{I}[i = a] \frac{(x_a - x_b)}{2p_a} + \mathbb{I}[i = b] \frac{(x_b - x_a)}{2p_b}$$

- It gives us a way to provide weight update rule in a concise form:

Weight update rule earlier 10: $w_a(t+1) \leftarrow w_a(t) \cdot e^{\frac{\gamma}{K} \frac{d}{2p_a}}$

11: $w_b(t+1) \leftarrow w_b(t) \cdot e^{-\frac{\gamma}{K} \frac{d}{2p_b}}$

Weight update rule using $\hat{c}_i(t)$ $\forall i \ w_i(t+1) = w_i(t) \cdot e^{\frac{\gamma}{K} \hat{c}_i(t)}$

Key element of the analysis

Lemma for expectation of $\hat{c}_i(t)$

$$\mathbb{E}[\hat{c}_i(t)|(a_1, b_1), \dots, (a_{t-1}, b_{t-1})] = x_i(t) - \mathbb{E}_{a \sim p(t)} x_a(t)$$

- The expectation of this estimator is the expected instantaneous regret of the algorithm against arm i .
- *i.e.* the difference between the gain of arm i and the expected gain according to algorithm's current state of knowledge $p(t)$.
- This is intuitively what we want from an estimator in a dueling bandit problem.

Sketch of proof

The general structure of the proof is similar to the proof of EXP3 [2] except the difference in expectation of the $\hat{c}_i(t)$ estimator.

Let $W_t = w_1(t) + w_2(t) + \dots + w_K(t)$.

$$\frac{W_{t+1}}{W_t} = \sum_{i=1}^K \frac{p_i(t) - \gamma/K}{1 - \gamma} e^{(\gamma/K)\hat{c}_i(t)} \quad (1)$$

As in EXP3, we simplify, take the logarithm and sum over t . We get for any j :

$$\sum_{t=1}^T \frac{\gamma}{K} \hat{c}_j(t) - \ln(K) \leq \frac{\gamma^2/K}{1 - \gamma} M_1 + \frac{(e - 2)\gamma^2/K}{1 - \gamma} M_2$$

Sketch of proof (continued)

By taking the expectation over the algorithm's randomization, we obtain for any j :

$$\begin{aligned} & \sum_{t=1}^T \frac{\gamma}{K} \mathbb{E}_{\sim p} \hat{c}_j(t) - \ln(K) \leq \\ & \frac{\gamma^2/K}{1-\gamma} \sum_{i=t}^T \mathbb{E}_{\sim p} M_1 + \frac{(e-2)\gamma^2/K}{1-\gamma} \sum_{i=t}^T \mathbb{E}_{\sim p} M_2 \end{aligned} \quad (2)$$

Sketch of proof (continued)

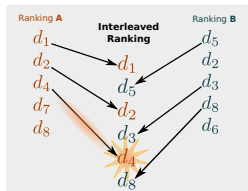
From Lemma 13, expectation of M_1 , expectation of M_2 , and by definition of \mathbb{G}_{max} , $\mathbb{E}\mathbb{G}_{alg}$, and $\mathbb{E}\mathbb{G}_{unif}$, the inequality (2) rewrites into:

$$\begin{aligned} \mathbb{G}_{max} - \mathbb{E}\mathbb{G}_{alg} - \frac{K \ln K}{\gamma} &\leq \frac{\gamma}{1-\gamma} (\mathbb{E}\mathbb{G}_{alg} - \mathbb{E}\mathbb{G}_{unif}) \\ &+ \frac{(e-2)\gamma}{2(1-\gamma)} (\mathbb{E}\mathbb{G}_{alg} + \mathbb{E}\mathbb{G}_{unif}) \end{aligned}$$

Assuming $\gamma \leq \frac{1}{2}$ we finally obtain:

$$\mathbb{G}_{max} - \mathbb{E}\mathbb{G}_{alg} \leq \frac{K \ln K}{\gamma} + \gamma (e\mathbb{E}\mathbb{G}_{alg} - (4-e)\mathbb{E}\mathbb{G}_{unif})$$

Experiments



- We used **interleaved filtering** on **real datasets** from information retrieval systems.
- We considered the following state of the art algorithms: **BTM** [6] (explore-then-exploit setting), **SAVAGE** [5], **RUCB** [7], and **SPARRING** coupled with **EXP3** [1] and **Random** as baseline.
- The experiments showed that **REX3** and especially its anytime version are competitive solutions for the dueling bandit problem.

Experiments

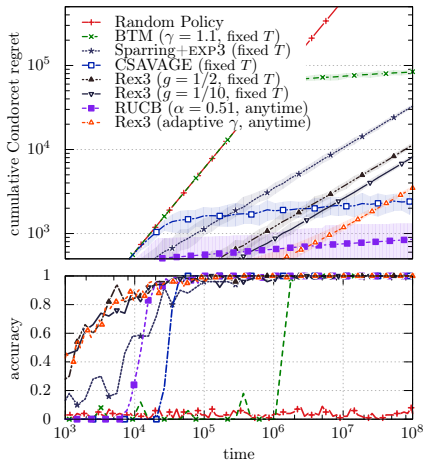


Figure 1: Average regret and accuracy plots on ARXIV dataset (6 rankers). Time and regret scales are logarithmic.

Experiments

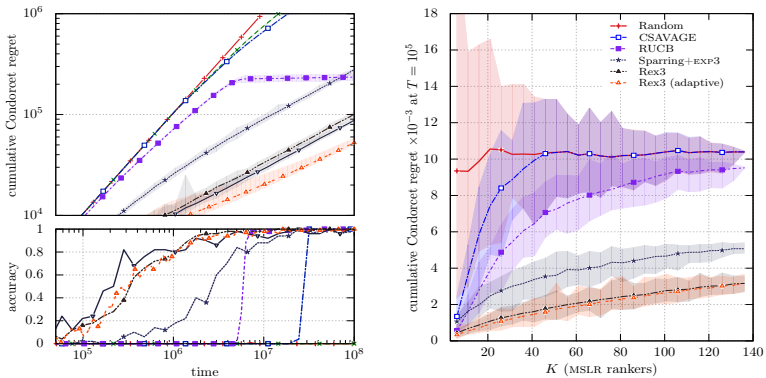


Figure 2: On the left: average regret and accuracy plots on MSLR30K with navigational queries ($K = 136$ rankers). On the right: same dataset, fixed $T = 10^5$ and $K = 4 - 136$. Colored areas show minimal and maximal values.

Simulations on non-stationary rewards

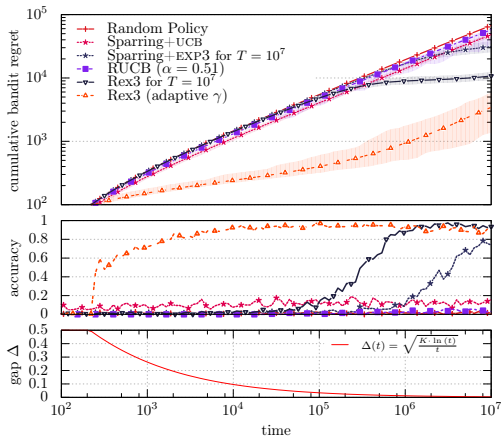


Figure 3: $K = 10$, gains from Bernoulli distributions. Best arm's gain is $1/2 + \Delta(t)$ with $\Delta(t) = \sqrt{K \cdot \log(t)/t}$. Others are stationary with a mean of $1/2$.

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